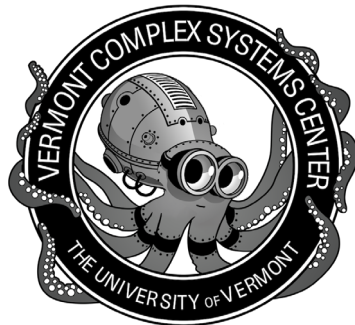


# The Role of Topology in Network Science

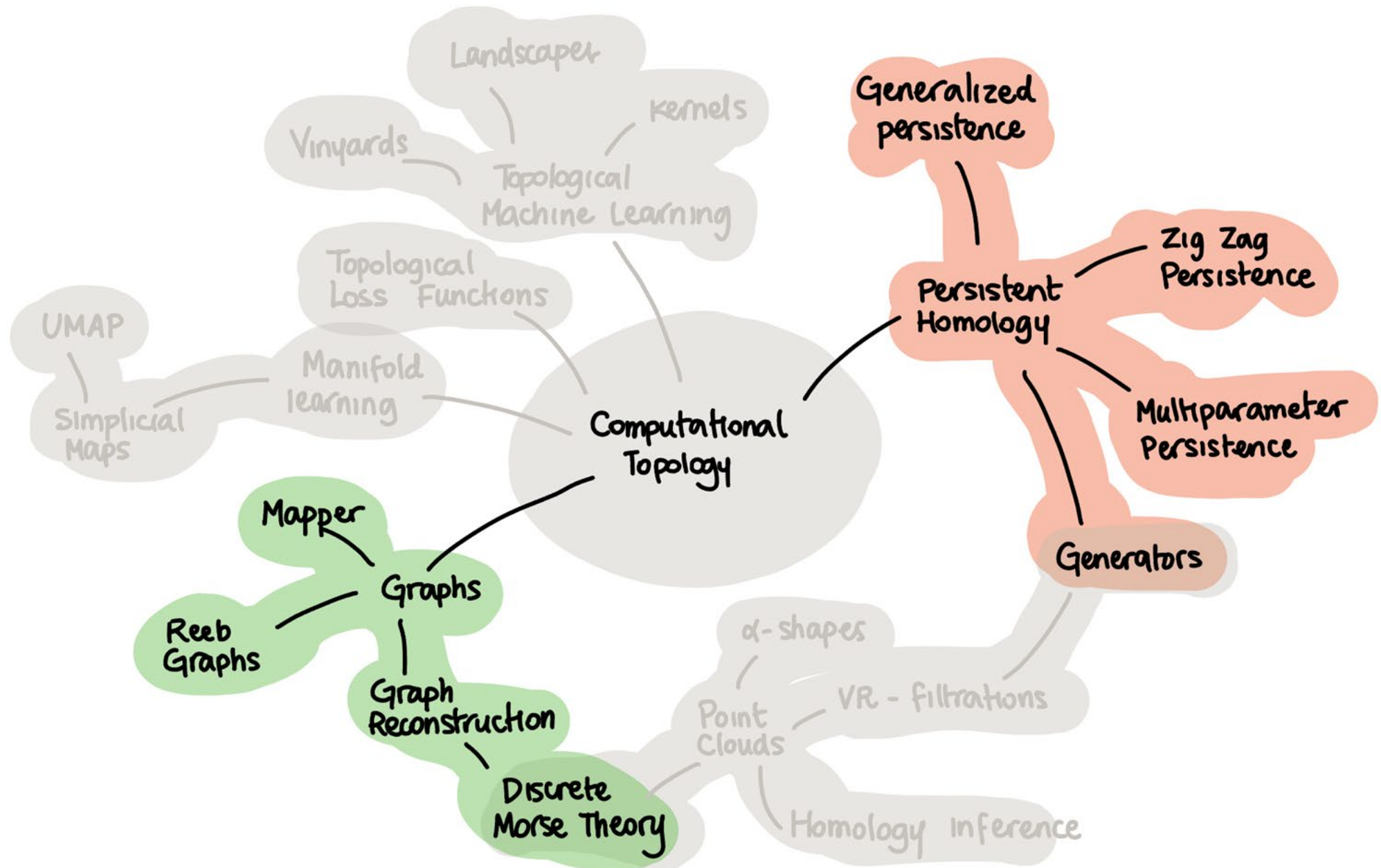
**Alice Patania**

Department of Mathematics and Statistics  
Vermont Complex Systems Center









# The Forking Path



The ancient forest looms before you, its gnarled trees stretching toward the sky like ancient guardians. Sunlight filters through the thick canopy, casting dappled shadows on the moss-covered ground. You have been following the faint trail for hours, and now you stand at a crossroads—a fork in the path that could lead you deeper into mystery or back to safety.

## Choice 1: The Whispering Grove

To your left, the path winds through a grove of silver-barked trees. You hear faint whispers carried on the breeze. Do you follow this path, drawn by curiosity and the promise of hidden knowledge? **(continue to Geometric Realizations)**

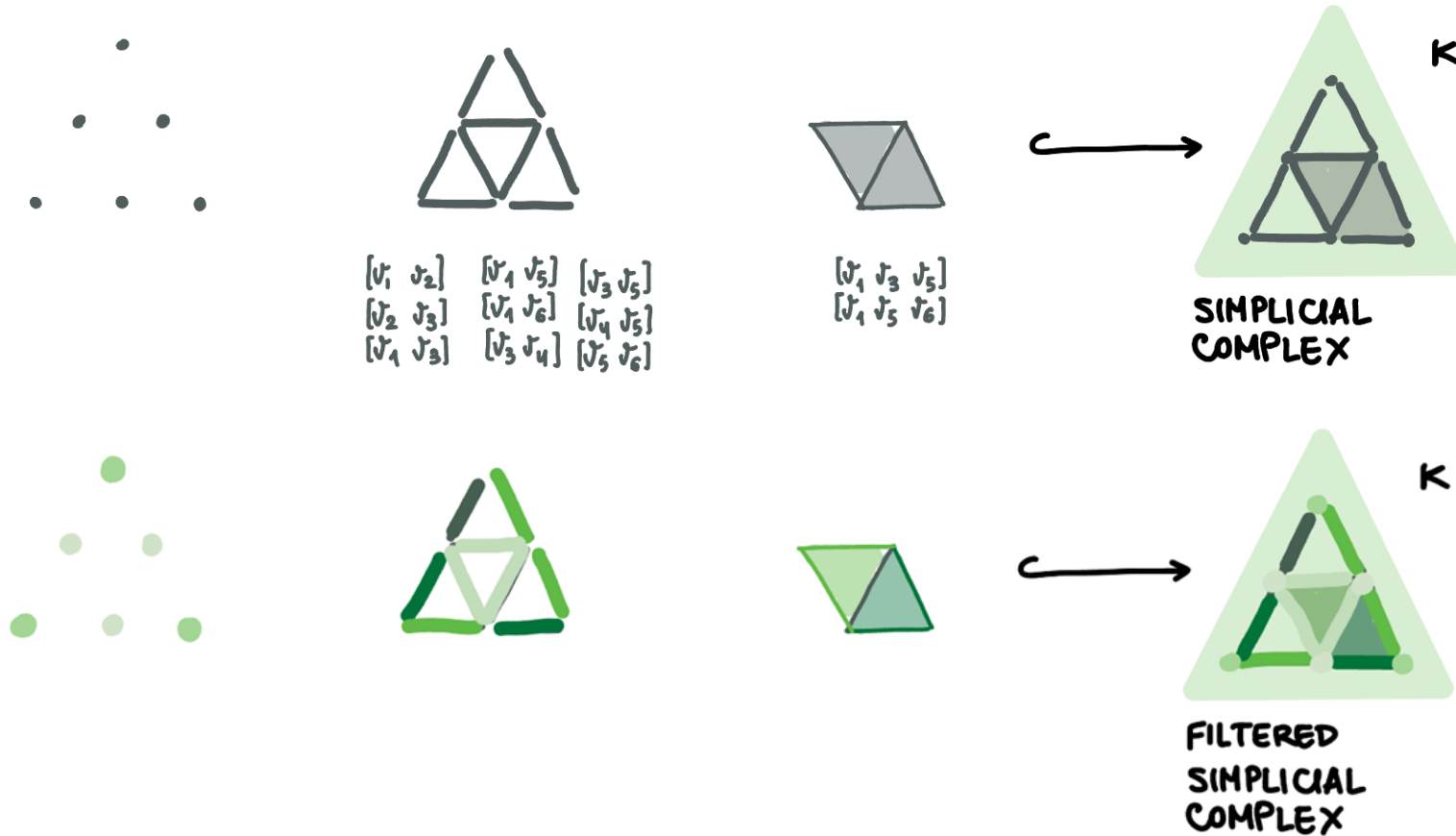
## Choice 2: The Forgotten Bridge

Straight ahead, a rickety wooden bridge spans a rushing stream. The water below is crystal clear, bringing forth memories of a forgotten past. Will you risk the bridge's creaking planks and venture toward fortune? **(continue to Topology 101)**

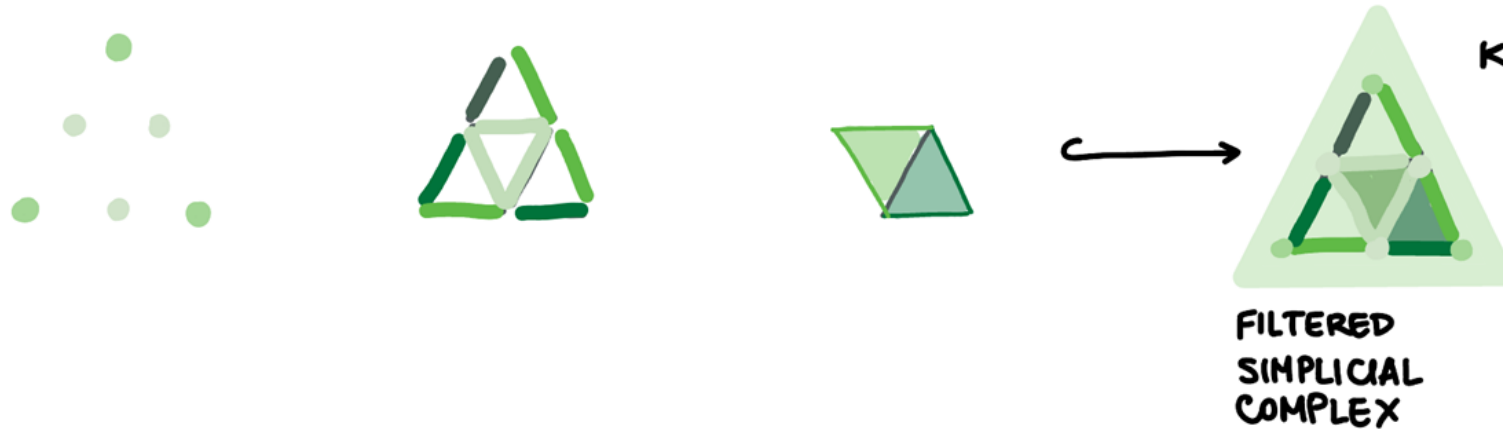
## Choice 3: You turn around and run home (turn to your phone or laptop)

Choose wisely, dear reader. Your fate awaits.

# At the beginning there were Higher-order relations



# Persistent Homology

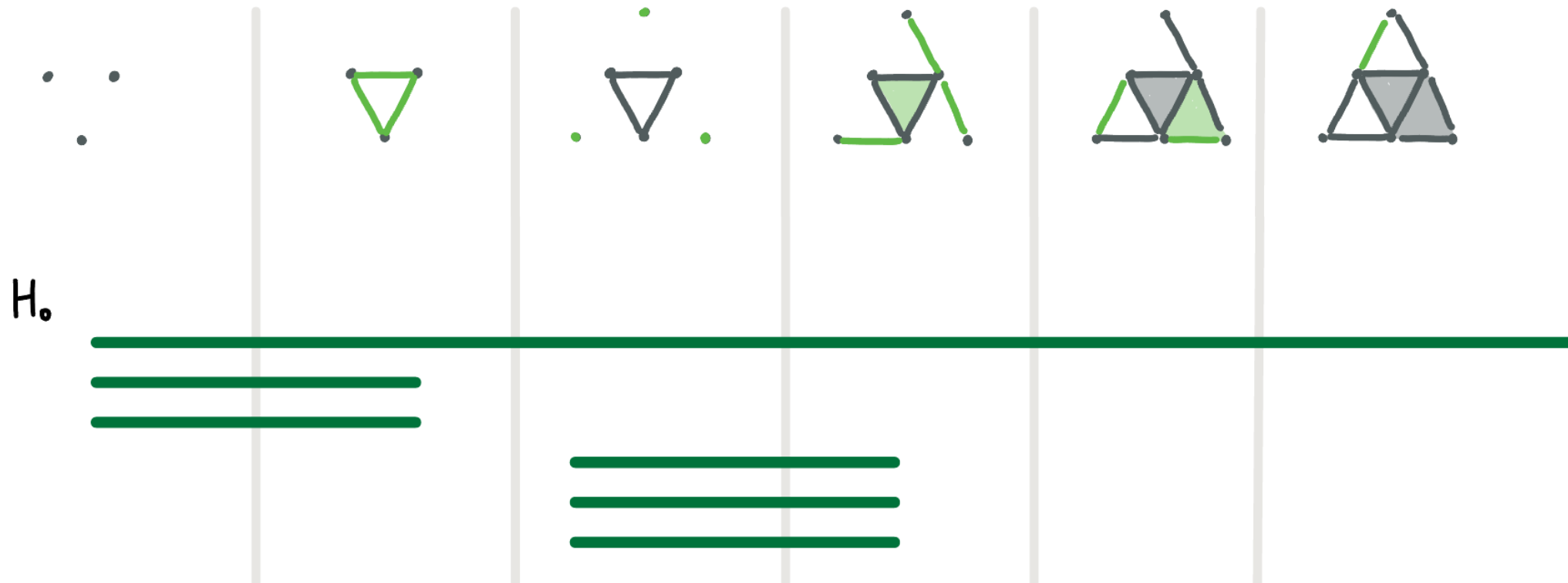


$$F_1 K \hookrightarrow F_2 K \hookrightarrow F_3 K \hookrightarrow F_4 K \hookrightarrow F_5 K \hookrightarrow F_6 K = K$$



# Persistent Homology

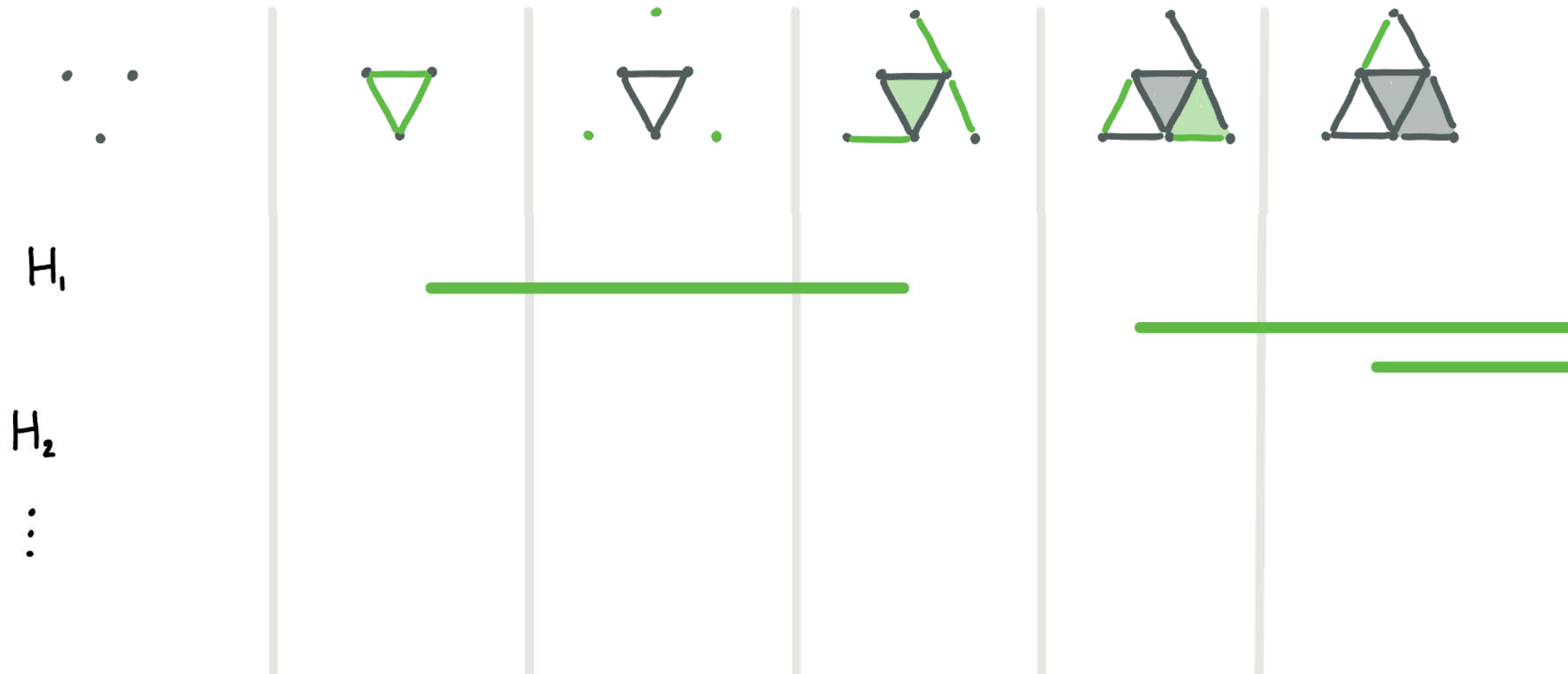
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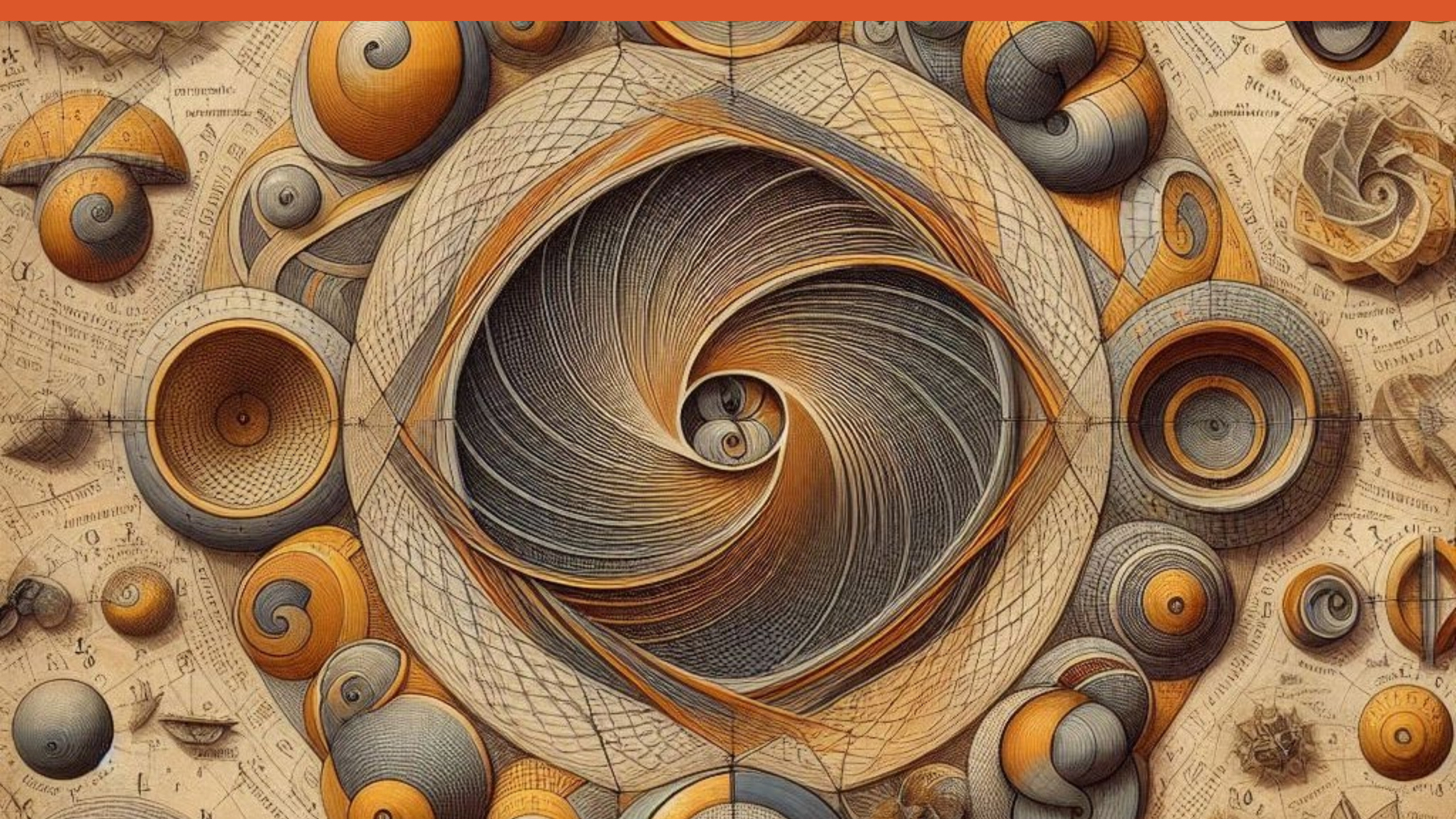




# Persistent Homology

$$F_1 K \hookrightarrow F_2 K \hookrightarrow F_3 K \hookrightarrow F_4 K \hookrightarrow F_5 K \hookrightarrow F_6 K = K$$

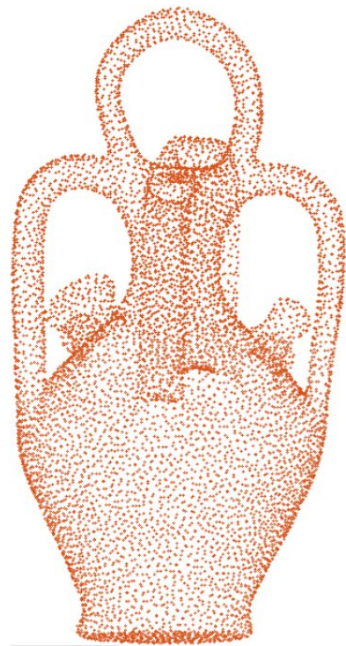




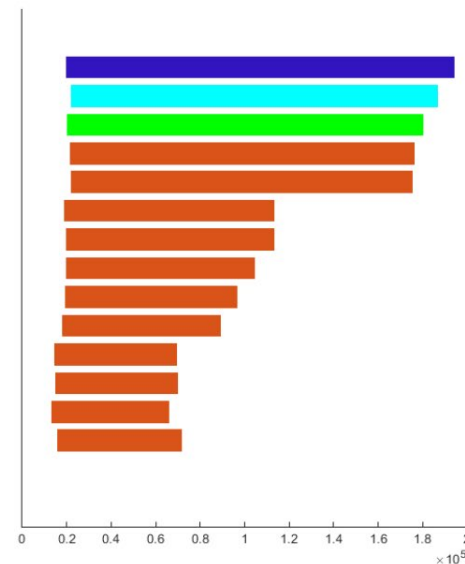
# Bringing Geometry back into Topology

Going beyond counting homological cycles

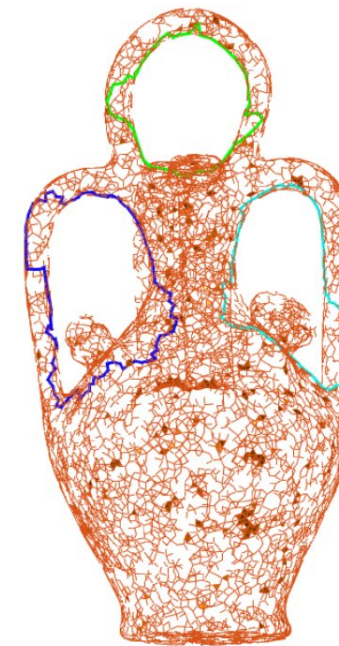
For the first decade of its life, TDA focused mostly on finding ways to count cycles and studying the stability under noise. In the last decade, focus has shifted to identifying a **optimal homology basis** [Eriksson Whittlesey 2005] and **minimal cycles** [Dey et al. 2018] [Guerra et al. 2021] [Li et al. 2021].



(a)



(b)



(c)

# Bringing Geometry back into Topology

Duality between persistent cycles and flow network cuts

Both these problems are NP-hard in general, but for **weak p-pseudomanifolds** \* optimal **p-1** homological cycles can be computed in polynomial time. [Chen and Freedman 2011]

A **flow network**  $(G, s, t, C)$  is an undirected graph  $G$ , two subsets of the vertices  $s, t \subset V(G)$  called *sources* and *sinks* respectively, and a *capacity* function  $C: E(G) \rightarrow [0, +\infty)$ .

A **cut** of  $G$  is a partition of the vertices  $S, T \subset V(G)$  so that  $s \subset S$ ,  $t \subset T$ , and  $S \cup T = V(G)$ .

The **capacity** of a cut  $(S, T)$  is the sum of the capacity of the edges between  $(S, T)$

$$C(S, T) = \sum_{e \in E(S, T)} C(e).$$

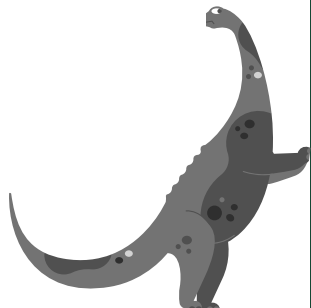
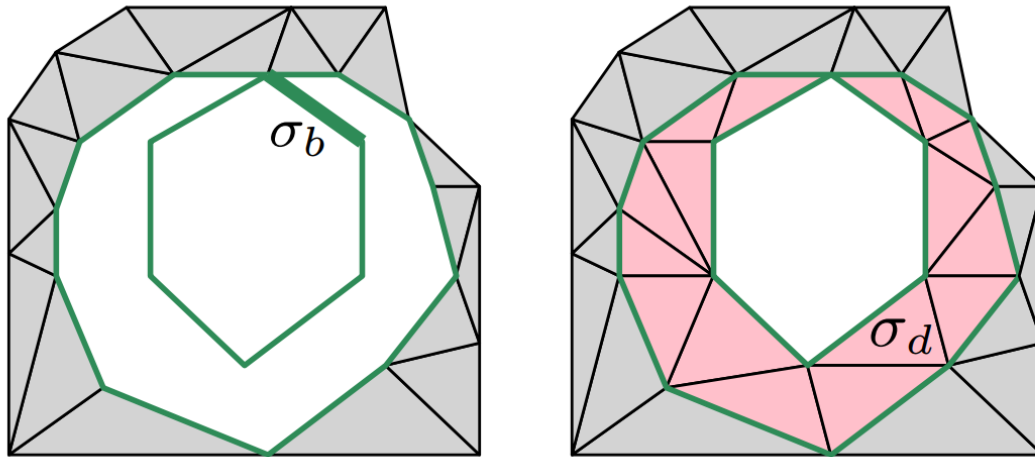
A **minimal cut** of  $(G, s, t, C)$  is a cut  $(S, T)$  with minimal capacity.

# Bringing Geometry back into Topology

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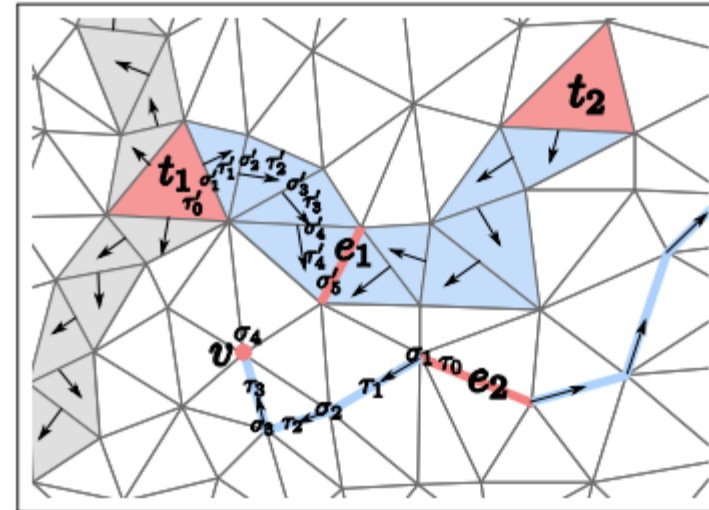
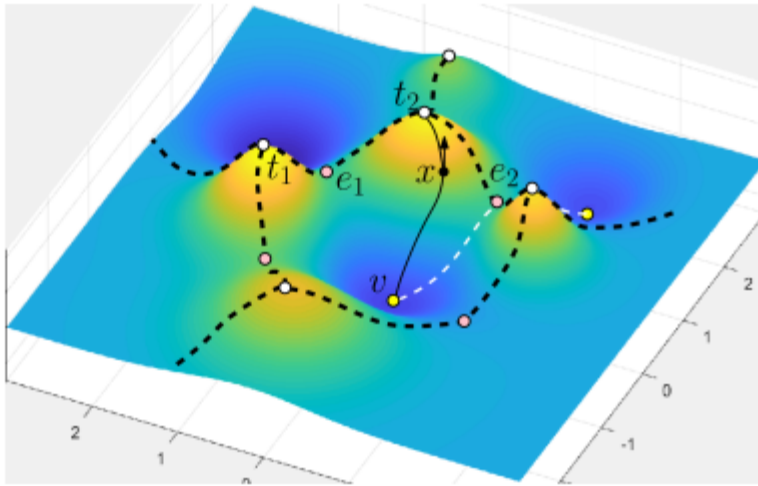
A **minimal cut** of  $(G, s, t, C)$  is a cut  $(S, T)$  with minimal capacity.



# The importance of a manifold

This intuition has been expanded to applications of Discrete Morse Theory to TDA.

**Morse Theory** is the branch of mathematics that connects the topology of smooth manifolds to the behavior of certain functions defined on those manifolds.

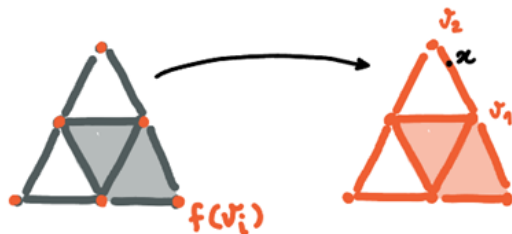


# The importance of a manifold

## Piecewise-Linear functions

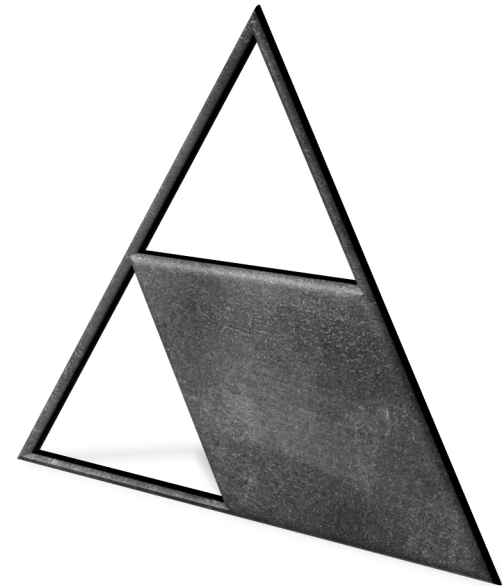
A **PL-function**  $\bar{f}: K \rightarrow \mathbb{R}$  is a real-valued function on the geometric realization of a simplicial complex determined by its restriction  $f: V \rightarrow \mathbb{R}$  to the vertex set  $V$  and linearly extending it within each simplex in  $K$ .

We have weights on the vertices



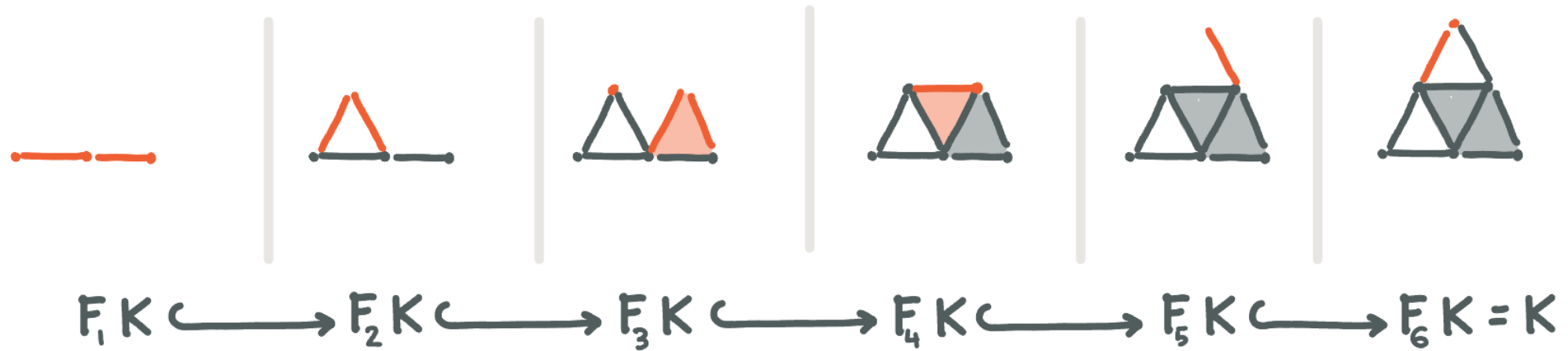
We extend them linearly to all  $K$

$$\bar{f}(x) = \sum_i^{k+1} a_i f(v_i) = \underbrace{\alpha_1 f(v_1) + \alpha_2 f(v_2)}_{\text{barycentric coordinate}}$$



# Piecewise-Linear functions

A **PL-function**  $\bar{f}: K \rightarrow \mathbb{R}$  defines a filtration on the geometric realization of  $K$



Closed star

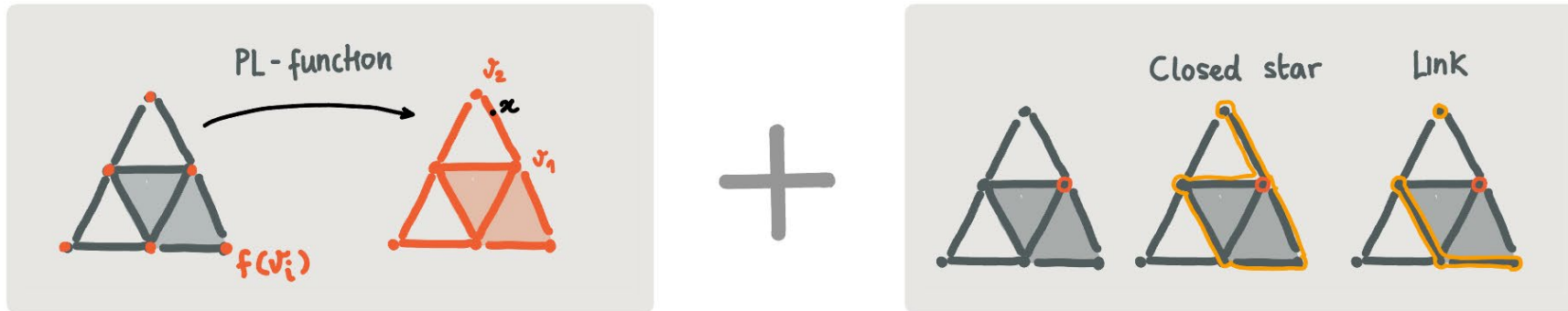
the simplicial complex induced by all the simplices that contain  $\heartsuit$

Link

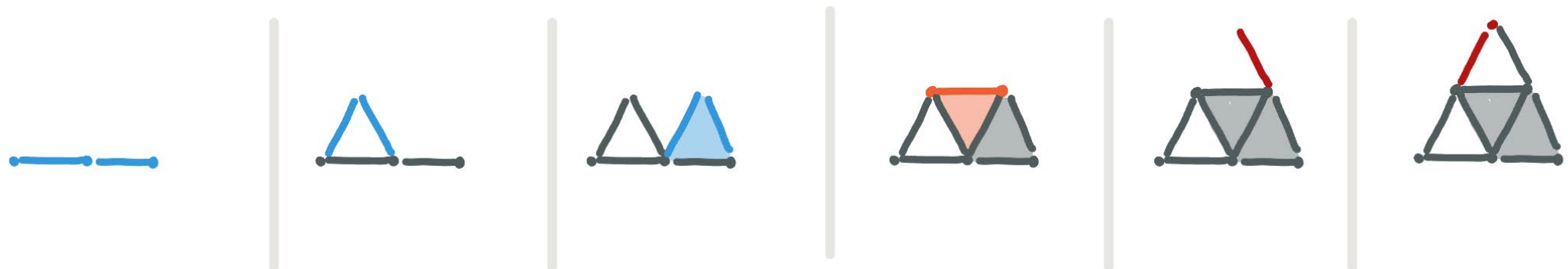
the simplices in the closed star that don't intersect  $\heartsuit$



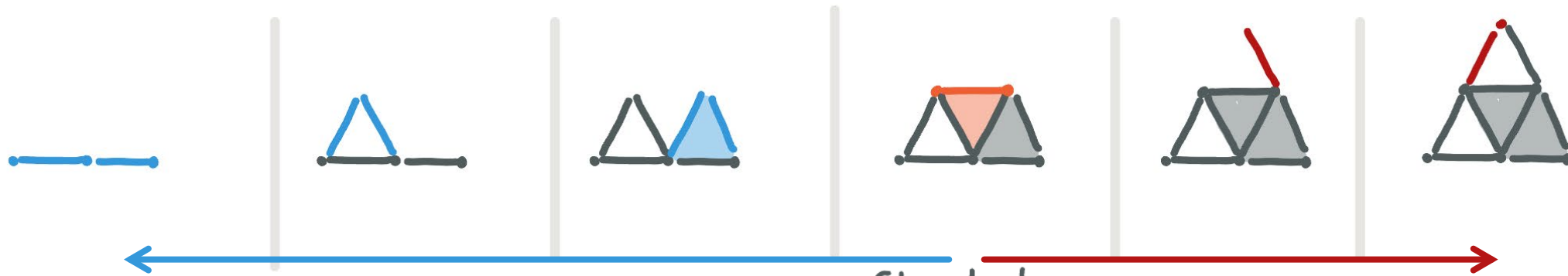
# Piecewise-Linear functions



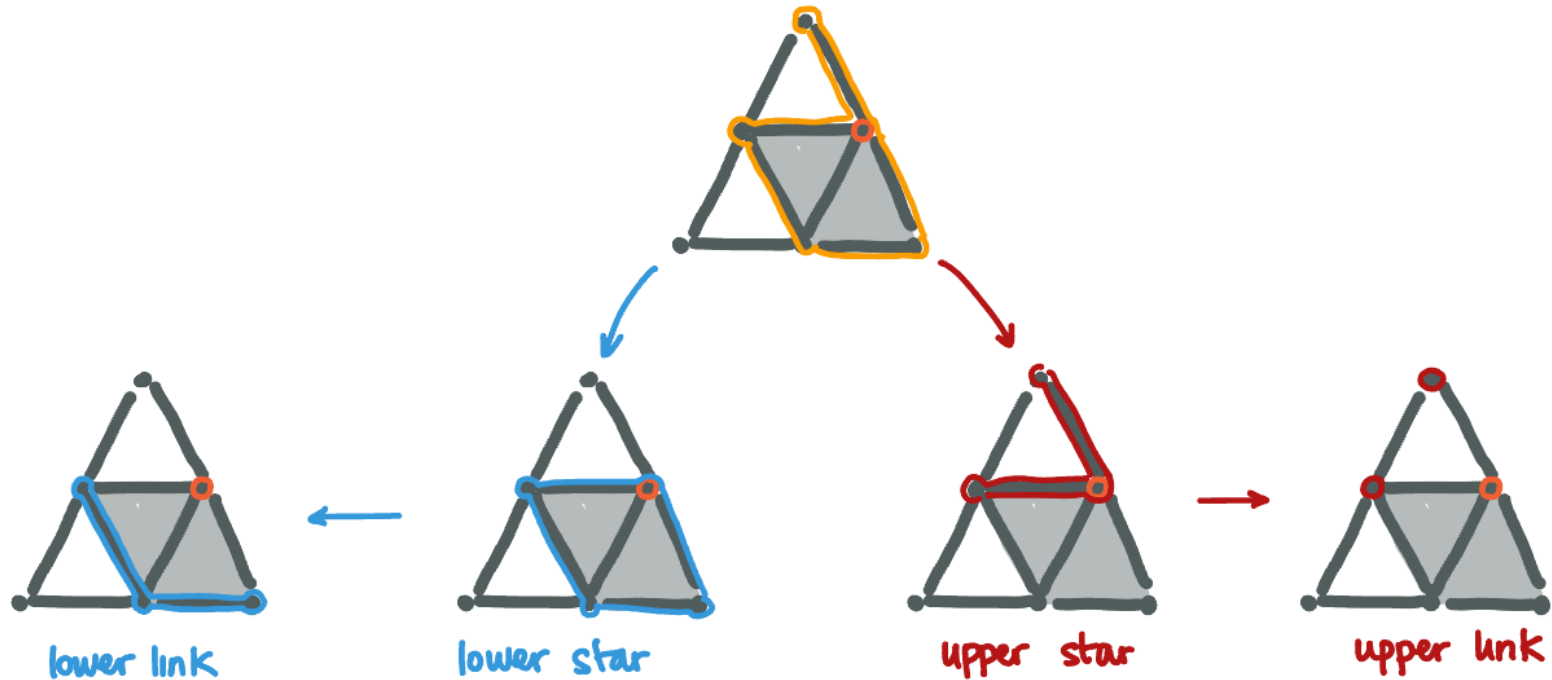
$$F_1 K \hookrightarrow F_2 K \hookrightarrow F_3 K \hookrightarrow F_4 K \hookrightarrow F_5 K \hookrightarrow F_6 K = K$$



# Piecewise-Linear functions



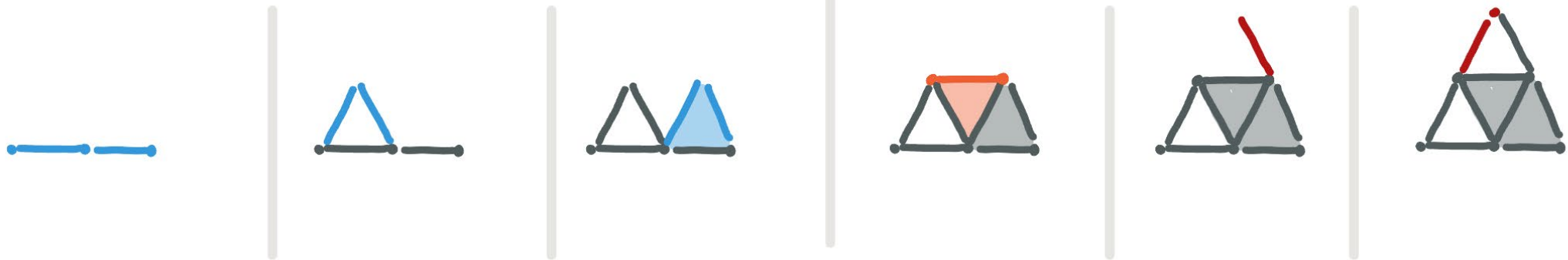
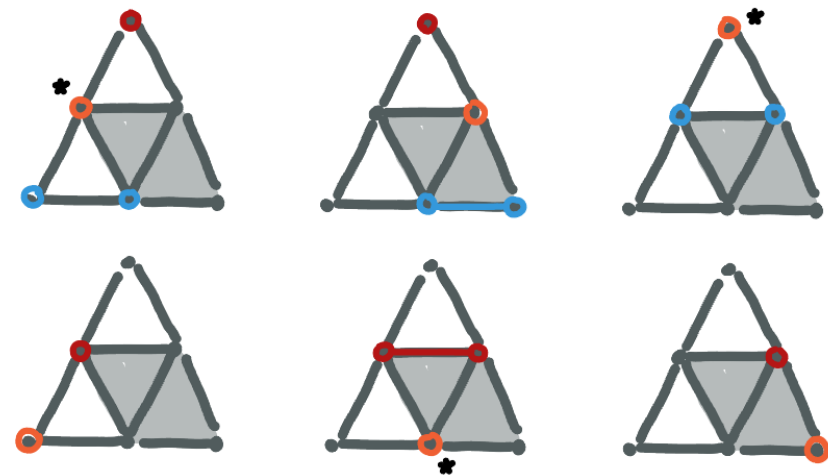
Closed star



# Critical Points

$$H_p \left( \begin{array}{c} \text{triangle with } v \text{ on lower link} \\ \text{lower link} \end{array} \right) \cong H_p \left( \begin{array}{c} \text{triangle with } v \text{ on upper link} \\ \text{upper link} \end{array} \right)$$

for all  $p > 0$  then  $v$  is a regular point of  $K$   
 otherwise  $v$  is a critical point of  $K$



# Why Should I care?

## PL-functions and Networks

Computing 0<sup>th</sup> persistent homology for a PL-function we retrieve the Kruskal's **Minimum Spanning Forest algorithm** for the graph underlying  $K$ .

This means that computing the persistent homology of a graph can be computed in  $O(n \log(n))$ . [Dey and Wang, 2022]

## Generalizations of Persistent Homology

There are 2 ways to generalize a filtration: the functions [Dey, Fan, Wang 2012] or the directions. [Carlsson and De Silva, 2010]\*

Graph Reconstruction  $F_1 K \hookrightarrow F_2 K \hookrightarrow F_3 K \hookrightarrow F_4 K \hookrightarrow F_5 K \hookrightarrow$

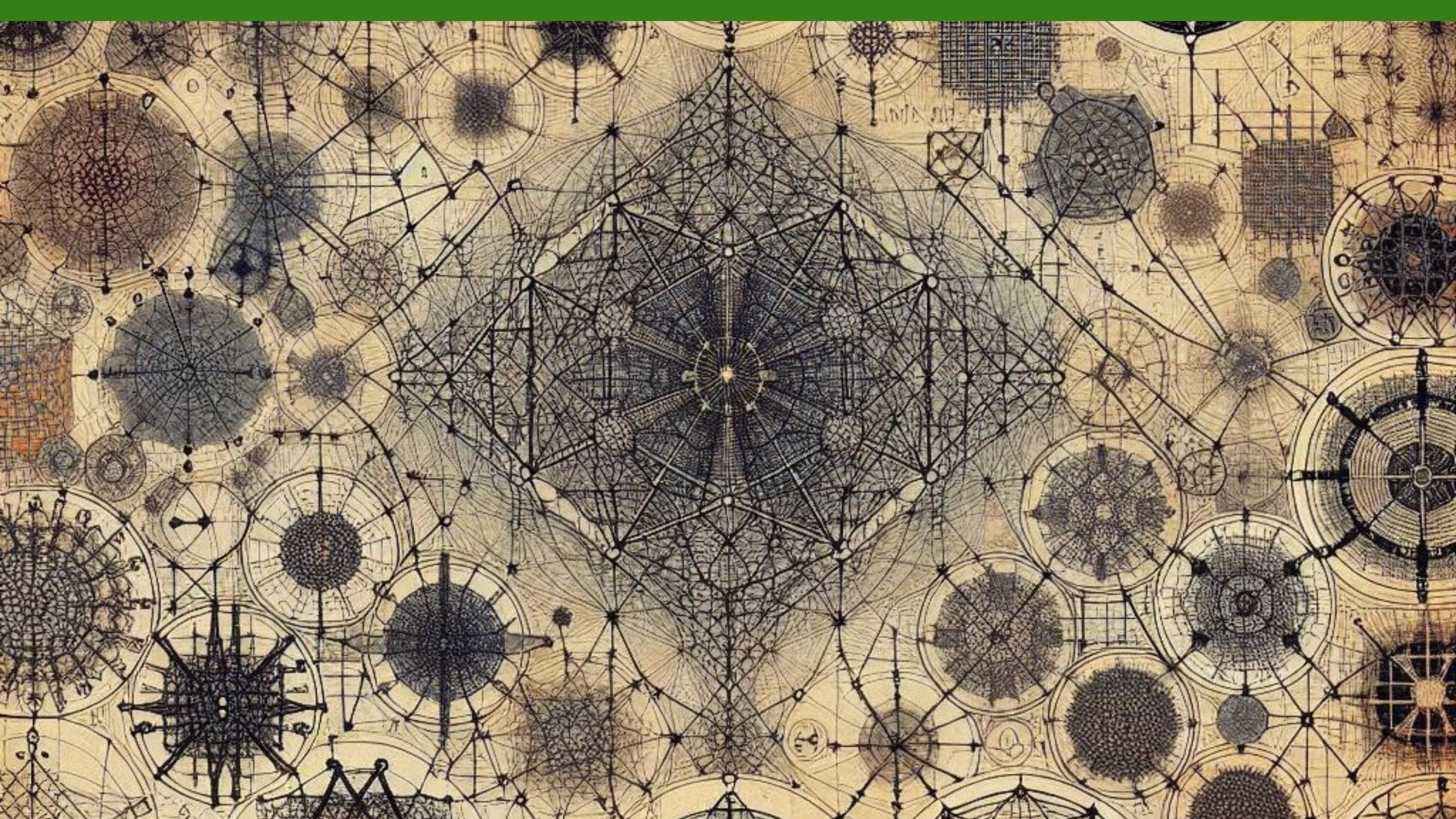
Suppose we have a hidden geometric graph  $G$  embedded in  $\mathbb{R}^n$  and a density PL-function  $\rho$  that concentrates around the hidden graph  $G$ . We can extend the Kruskal algorithm to reconstruct  $G$  from  $\rho$ . [Wang et al., 2015]\*\*

zig-zag  
filtrations



simplicial  
towers





# Mapper

MAPPER



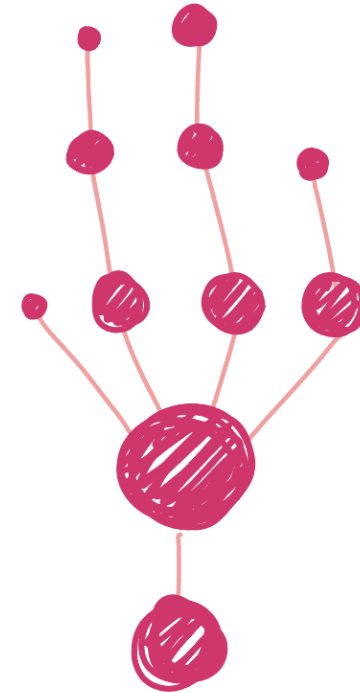
Point cloud



Real function



bns

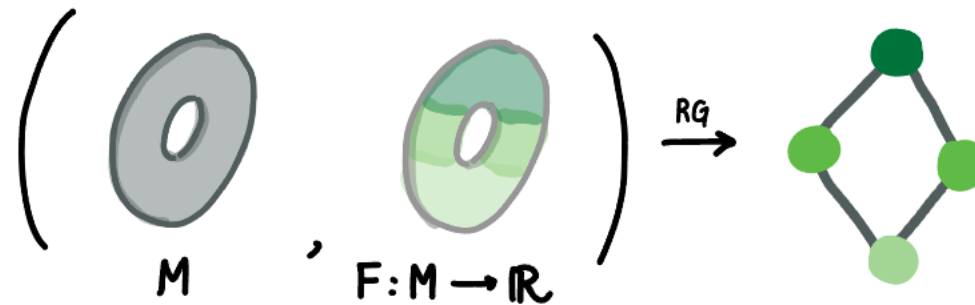


graph

[Dey Memoli Wang (2017) *Topological analysis of nerves, Reeb spaces, mappers, and multiscale mappers*]

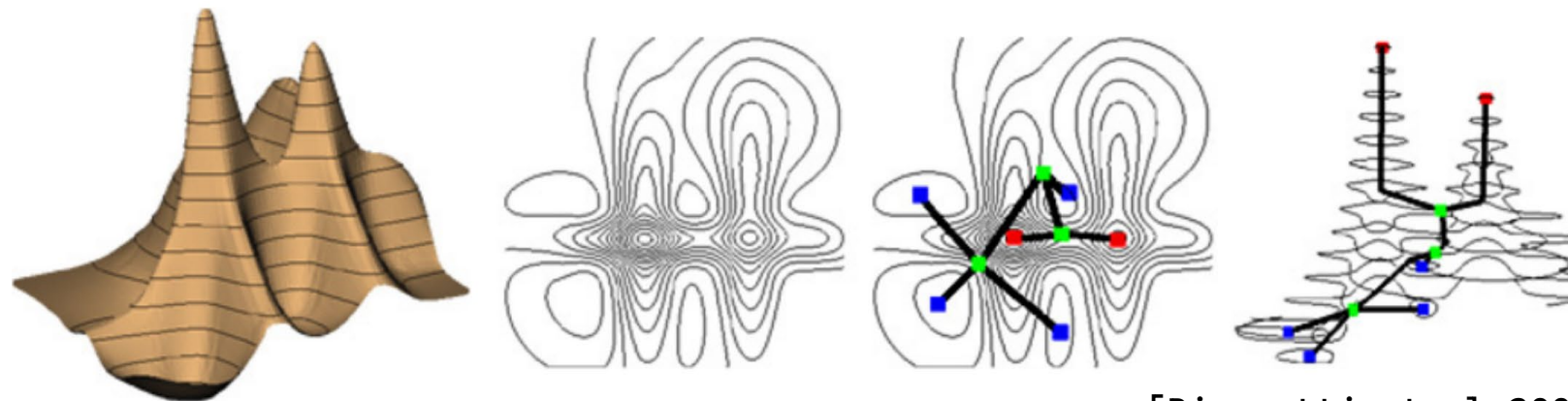
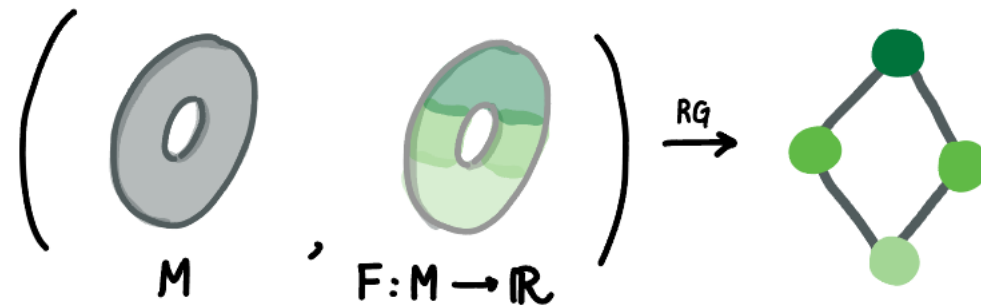
# Reeb Graphs (contour trees in cluster analysis)

Mapper is an algorithm related to Reeb Graphs. These are combinatorial objects used to study a real-valued function defined on a manifold.



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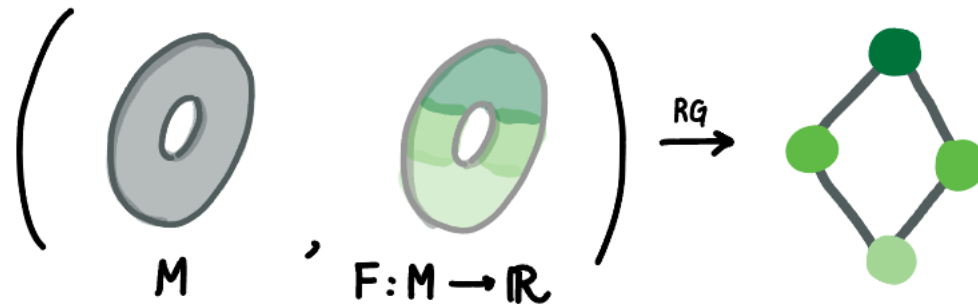


[Biassotti et al 2008]



# Reeb Graphs (contour trees in cluster analysis)

Mapper is an algorithm related to Reeb Graphs. These are combinatorial objects used to study a real-valued function defined on a manifold.

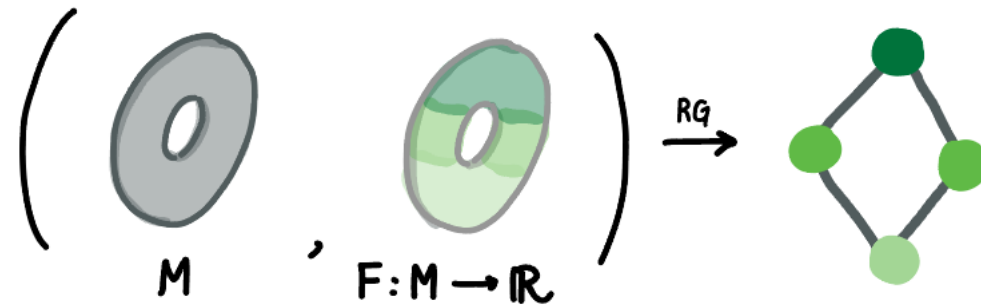


A lot of the research on applications of Reeb Graphs is around finding ways to compare them with one another.

- Interleaving distance [Morozov et al. 2013]
- Functional Distortion distance [Bauer et al. 2014]
- Edit distance [Di Fabio, Landi 2016]
- Bottleneck distance [Cohen et al. 2006]

# Reeb Graphs (contour trees in cluster analysis)

Mapper is an algorithm related to Reeb Graphs. These are combinatorial objects used to study a real-valued function defined on a manifold.



To compute the Reeb graph of a **simplicial complex**, the best approach is to use PL-functions in  $O(m \log m)$ . [Parsa 2013]

Why should we care?

As a community we are starting to look at Mapper to study networks

Maybe looking at how its continuous counterpart has been used might help us.

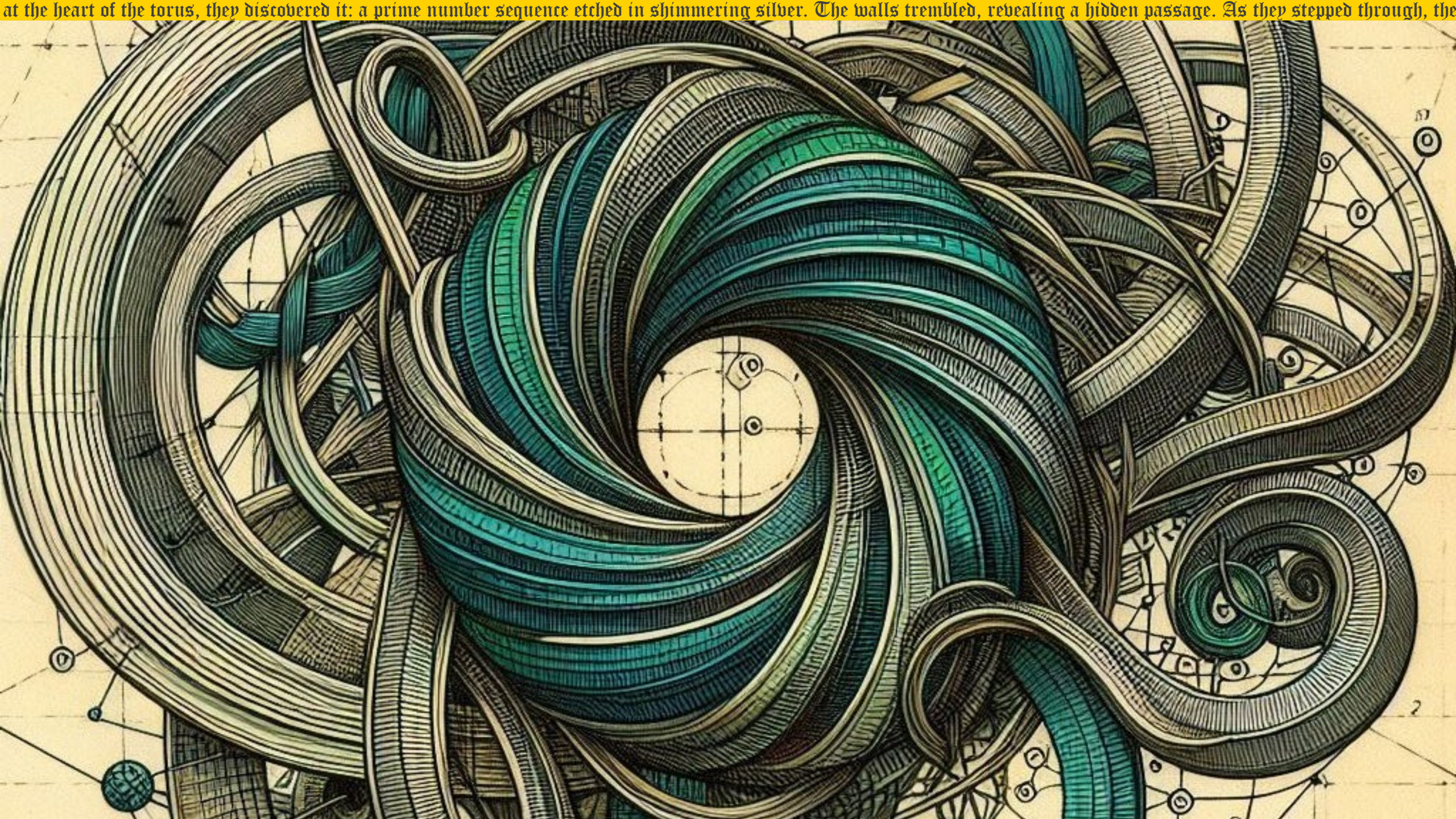
[Hajij et al 2018]

[Rosen et al 2018]

[Bodnar et al 2020]

[Patania et al 2023ish]

at the heart of the torus, they discovered it: a prime number sequence etched in shimmering silver. The walls trembled, revealing a hidden passage. As they stepped through, the

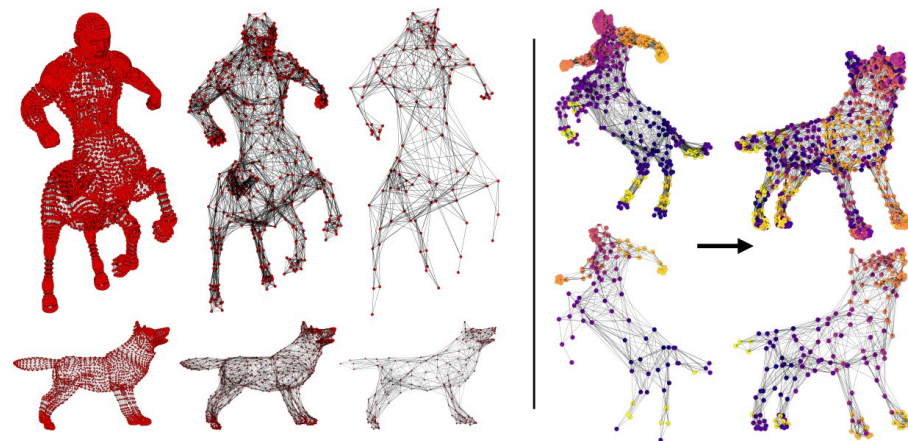


# Esoteric Mapper

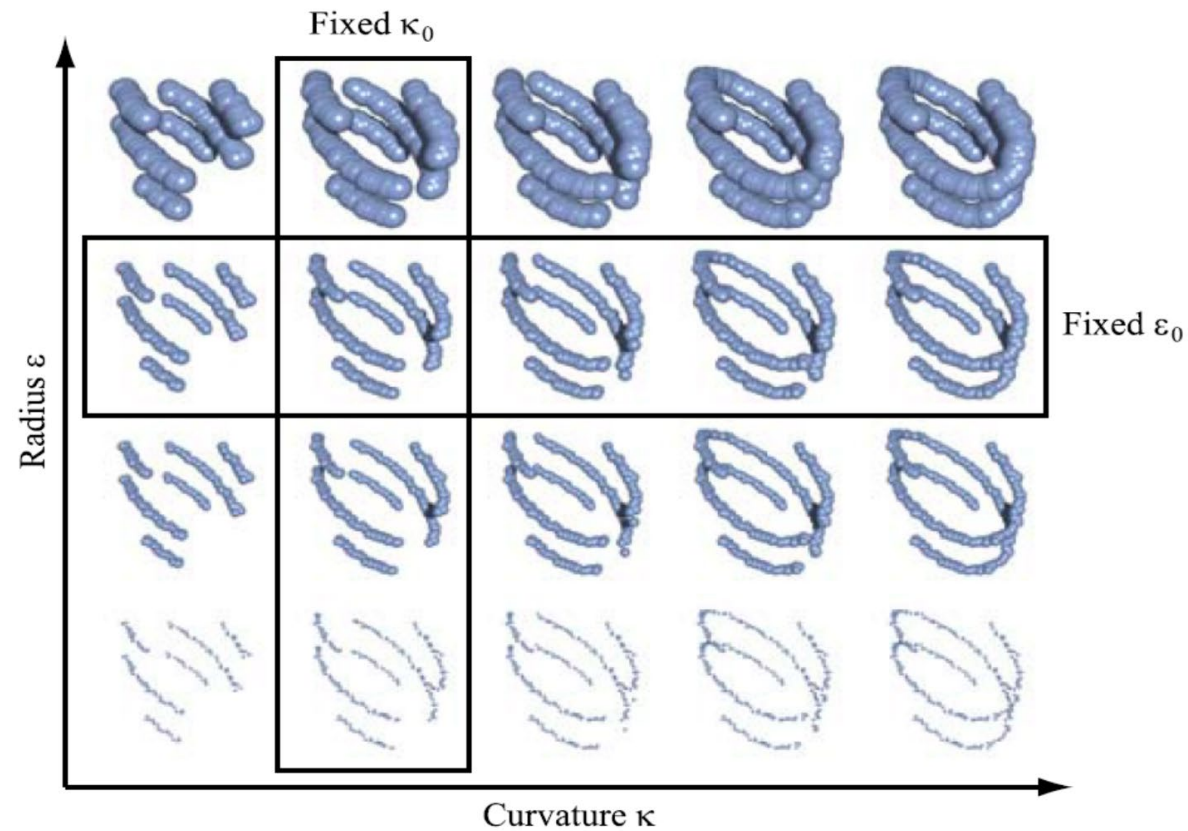
We have always seen Mapper as a graph, but it can be defined as the skeleton of more general functions and covers of a point cloud, turning it into a proper simplicial complex. [Singh et al. 2007]

In 2016, Dey et al. created the mythological monster that is **multiscale mapper** putting together ideas from generalized persistence with that of combinatorial mapper. Unaware of its existence, Chowdhury et al in 2021 used it with co-optimal transport to match combinatorial structures generated at different resolutions.

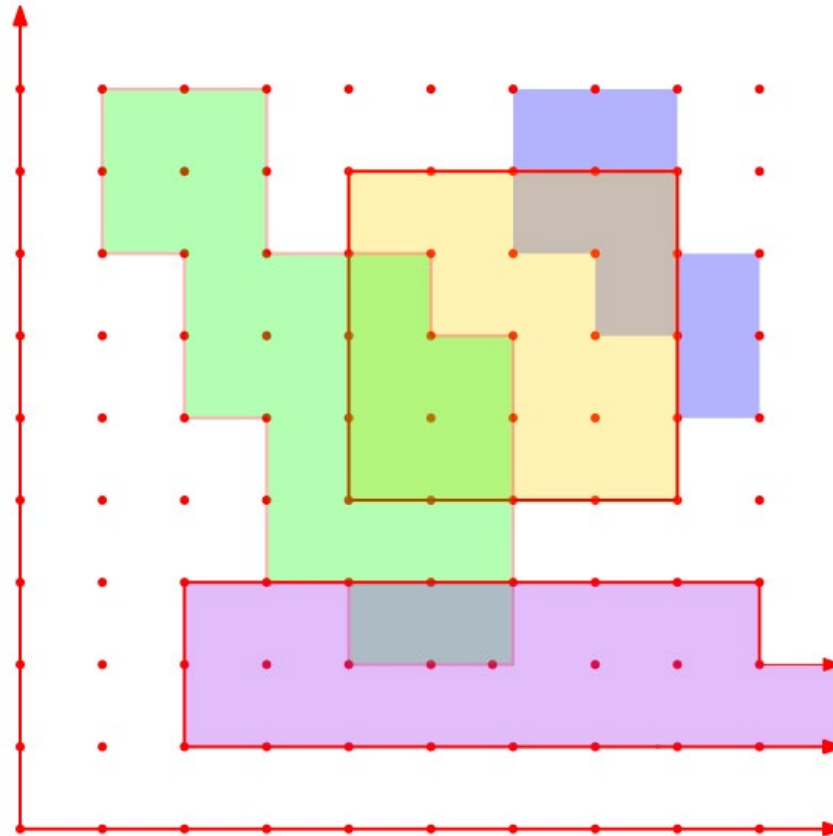
Multi-resolution Hypergraph Matching



# Multiparameter persistence – adding dimensions



# Multiparameter persistence – adding one dimension



# Multiparameter persistence – adding one dimension

$$\text{Rk} \begin{pmatrix} \text{k} & \xrightarrow{\text{id}} & \text{k} & \longrightarrow & 0 \\ \text{id} & & \uparrow & & \uparrow \\ & & \text{k} & \xrightarrow{[1\ 0]} & \text{k} \\ & & \downarrow & & \downarrow \\ \text{k} & \xrightarrow{[0]} & \text{k}^2 & \xrightarrow{[1\ 1]} & \text{k} \\ & & \downarrow & & \downarrow \\ 0 & \longrightarrow & \text{k} & \xrightarrow{\text{id}} & \text{k} \end{pmatrix} = \text{Rk} \left( \begin{array}{c} \text{Diagram 1} \oplus \text{Diagram 2} \oplus \text{Diagram 3} \end{array} \right) - \text{Rk} \left( \begin{array}{c} \text{Diagram 4} \end{array} \right)$$

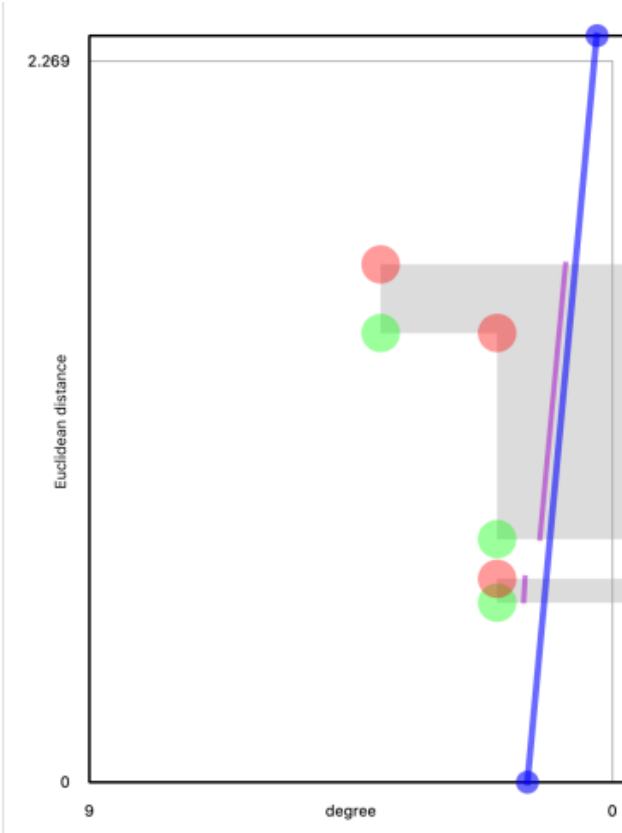
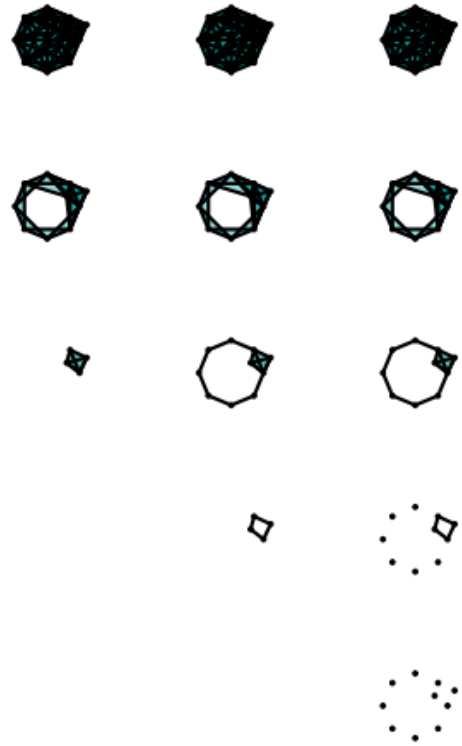
The diagrammatic equation shows the rank of a complex of vector spaces over a field  $\mathbb{k}$ . The complex is represented by a grid of nodes and arrows. The rank is calculated as the rank of the direct sum of three diagrams (shaded blue) minus the rank of one diagram (shaded red). The blue diagrams represent the rank of the subcomplexes, and the red diagram represents the rank of the quotient complex.

$$\text{Rk} \begin{pmatrix} \text{k} & \xrightarrow{\text{id}} & \text{k} & \longrightarrow & 0 \\ \text{id} & & \uparrow & & \uparrow \\ & & \text{k} & \xrightarrow{[1\ 0]} & \text{k} \\ & & \downarrow & & \downarrow \\ \text{k} & \xrightarrow{[0]} & \text{k}^2 & \xrightarrow{[1\ 1]} & \text{k} \\ & & \downarrow & & \downarrow \\ 0 & \longrightarrow & \text{k} & \xrightarrow{\text{id}} & \text{k} \end{pmatrix} = \text{Rk} \left( \begin{array}{c} \text{Diagram 1} \oplus \text{Diagram 2} \oplus \text{Diagram 3} \oplus \text{Diagram 4} \end{array} \right) - \text{Rk} \left( \begin{array}{c} \text{Diagram 5} \oplus \text{Diagram 6} \end{array} \right)$$

The diagrammatic equation shows the rank of the same complex as above, but with a different decomposition. The rank is calculated as the rank of the direct sum of four diagrams (shaded blue) minus the rank of two diagrams (shaded red). The blue diagrams represent the rank of the subcomplexes, and the red diagrams represent the rank of the quotient complex.

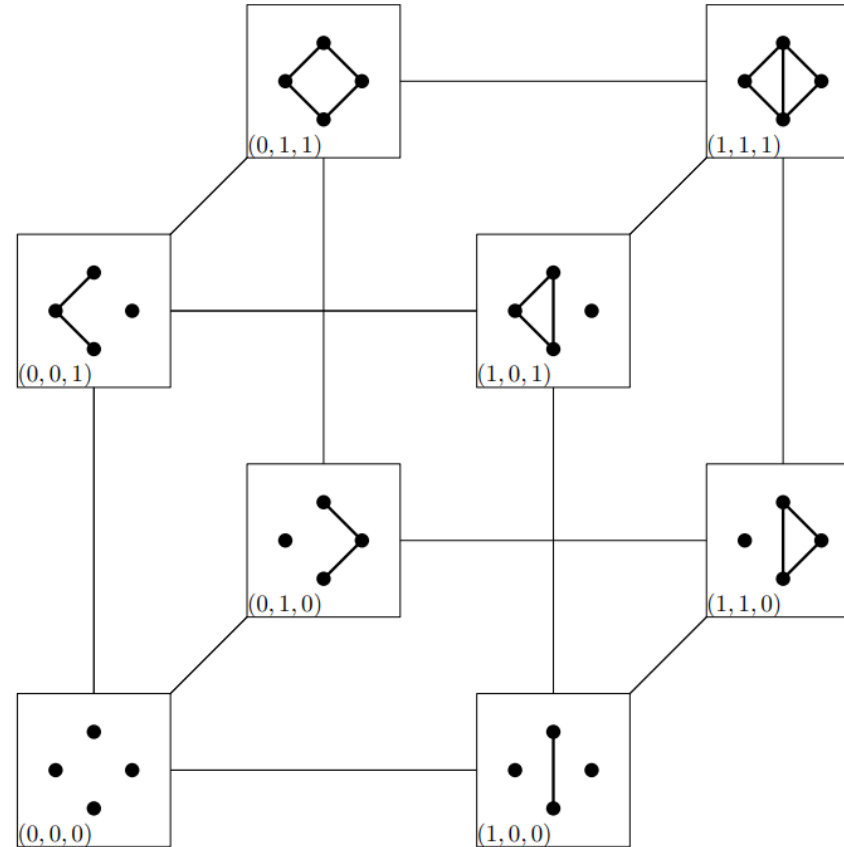


# Multiparameter persistence – adding one dimension





# Multiparameter persistence – adding two dimensions

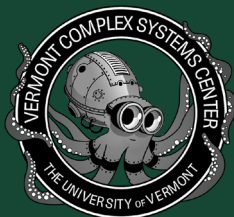
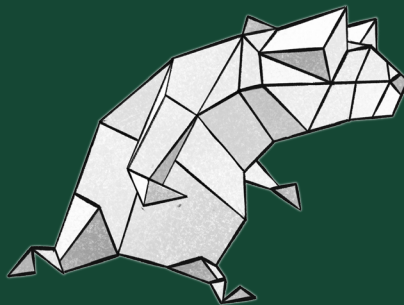
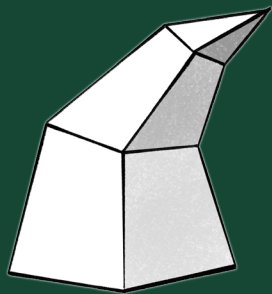
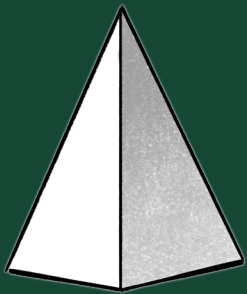


What can you use it for?

# What can you use it for?



# Thank You



The End